

Network Economics and Security Engineering

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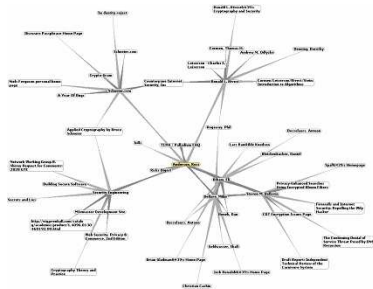
Outline

- 1 Relevant network properties
- 2 Example Applications
- 3 Conclusions

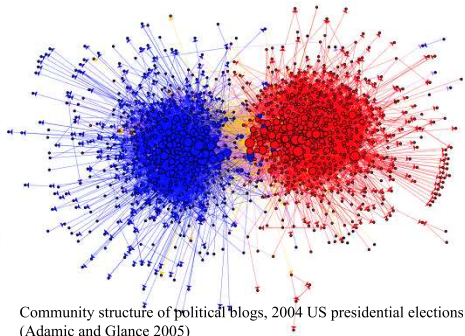
Motivation

- Many computing applications are network-based
 - World-wide web
 - Internet backbone
 - Peer-to-peer networks
 - Wireless sensor networks
 - Social networks
- Network externalities matter: the decisions of others impact a user's best response
- Interactions on these networks can be modeled as repeated games with evolving strategies
- Network properties influence dominant strategy outcomes

How do we represent a network?



<http://www.cl.cam.ac.uk/~rja14/>

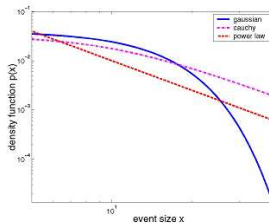
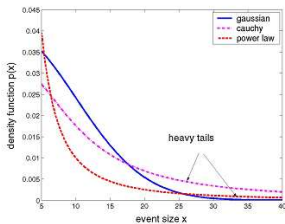


Community structure of political blogs, 2004 US presidential elections (Adamic and Glance 2005)

Relevant network properties

- Network topology
- Network dynamics
- Adversarial model
 - Different attacker goals
 - Different attacker capabilities

Network topology



- Fully-connected graph
- Lattice
- Random graph (Erdős-Rényi)
- Geometric random graph
- Scale-free degree distribution
- Small-world topology

Network dynamics

- Node mobility
 - Lessens likelihood of repeated interaction
 - Allows malicious nodes to maximize attack
- Churn
 - Lessens likelihood of repeated interaction
 - Makes punishment by exclusion difficult
 - Makes Sybil attacks likely
- Intermittent connectivity
 - Makes fair resource contribution difficult to establish
- Each of the above dynamics creates an informational asymmetry

Attacker goals

- Network partition
 - Good strategy for a communications network, maybe not for a file-sharing network with built-in redundancy
- Disrupt operations of 'normal' protocols (e.g., message routing)
- Avoid detection and punishment
- Maximize eavesdropping capability

Attacker capabilities

- Global knowledge
 - Powerful adversary can identify central nodes
- Local knowledge
 - Random walk to infer network topology

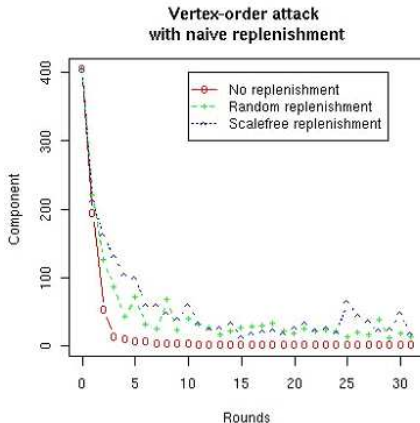
Topology of covert conflict

- Scale-free network
- No application-level network dynamics (mobility, churn, etc.)
- Attacker goal: network partition
- Defender goal: maximize connectivity
- Attacker has global knowledge of network topology, defender has local knowledge
- Goal: study interaction between dynamic attack and defense strategy

Attack and defense mechanisms

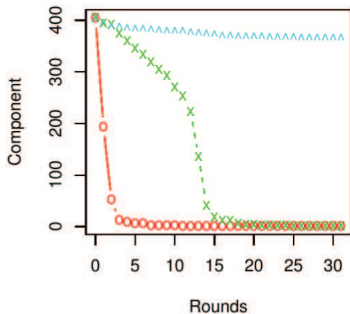
- Attack mechanisms
 - Remove nodes with high degree
 - Remove nodes with high betweenness centrality
- Defense mechanisms
 - Naive replenishment
 - Localized rings
 - Localized cliques

Attack under naive replenishment



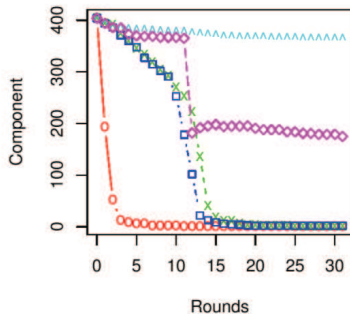
Attack and ring/cliq defenses

Vertex-order attack with Rings and Cliques



- vertex order attack, No adaptation
- +— vertex order attack, Rings
- △— vertex order attack, Cliques

Vertex-order and Centrality attack with Rings and Cliques



- vertex order attack, No adaptation
- +— vertex order attack, Rings
- △— vertex order attack, Cliques
- centrality attack, Rings
- ◇— centrality attack, Cliques

- We would like to apply a similar repeated game simulation framework to other network applications
 - Vary attacker models and goals
 - Vary network topology and dynamics to study effect on security mechanisms
 - Test viability of security mechanisms by varying strategies
- Promising applications
 - Communications surveillance by a limited adversary(Danezis and Wittneben, WEIS 2006)
 - Punishment mechanisms in decentralized computer networks

Punishment mechanisms in decentralized computer networks

- When devices misbehave, often there is no central authority available to identify and punish malicious behavior
- Solution: collective-decision mechanism
 - Reputation system
 - Blackballing with threshold voting
- Attacker goals
 - Avoid punishment while misbehaving
 - Abuse strategy to disconnect the network
- We have explored this space, and have proposed alternative mechanisms

Alternative mechanisms for addressing misbehavior

- Blackballing
 - Nodes cast accusatory votes upon observing misbehavior; once enough votes are cast against a node, it is removed
- Reelection
 - Devices cast positive votes affirming their friends; strangers only interact when they can demonstrate having a sufficient number of friends
- Suicide
 - Nodes unilaterally decide when to remove a malicious node, but must sacrifice itself to demonstrate its sincerity

Open questions in the strategy space

- For blackballing and reelection, nodes can individually set thresholds according to their risk attitude
- Network topology and dynamics determine which strategy works best
 - Scale-free degree distribution makes high-degree nodes immune to thresholds and low-degree nodes susceptible
 - Likewise, suicide can be used to target high-value nodes
- Which strategy, if any, will dominate?

Potential game frameworks

- Hiding attacker
 - Initially, half of nodes are assigned to each strategy
 - Small fraction of nodes set as malicious (attacker goal: avoid punishment)
 - Each round:
 - 1 Attack (some nodes misbehave)
 - 2 Defend (implement strategy)
 - 3 Adapt strategy if node identifies an unpunished neighbor
- Active attacker
 - Initially, half of nodes are assigned to each strategy
 - Small fraction of nodes set as malicious (attacker goal: remove as many honest nodes as possible)
 - Each round:
 - 1 Attack (malicious nodes falsely accuse honest nodes)
 - 2 Defend (honest nodes try to punish attackers)
 - 3 Probabilistically change strategy if node identifies unpunished neighbors

Conclusions

- The structure and dynamics of networks can vary greatly
- It is not well understood how differences in network composition impact secure operation
- Simulations using a repeated game framework looks promising
- Much more work to be done!

More...

- <http://www.cl.cam.ac.uk/~twm29/>